

# Combining Static and Dynamic Data in Code Visualization

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## ABSTRACT

The task of developing, tuning, and debugging compiler optimizations is a difficult one which can be facilitated by software visualization. There are many characteristics of the code which must be considered when studying the kinds of optimizations which can be performed. Both static data collected at compile-time and dynamic runtime data can reveal opportunities for optimization and affect code transformations. In order to expose the behavior of such complex systems, visualizations should include as much information as possible and accommodate the different sources from which this information is acquired.

This paper presents a visualization framework designed to address these issues. The framework is based on a new, extensible language called *JIL* which provides a common format for encapsulating intermediate representations and associating them with compile-time and runtime data. We present new contributions which extend existing compiler and profiling frameworks, allowing them to export the intermediate languages, analysis results, and code metadata they collect as *JIL* documents. Visualization interfaces can then combine the *JIL* data from separate tools, exposing both static and dynamic characteristics of the underlying code. We present such an interface in the form of a new web-based visualizer, allowing *JIL* documents to be visualized online in a portable, customizable interface.

## Categories and Subject Descriptors

D.3.4 [Programming Languages]: Processors—*compilers, optimization*

## General Terms

Performance, Design, Experimentation, Languages.

## Keywords

Combining static and dynamic data, visualization, intermediate languages, profiling, software understanding.

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## 1. INTRODUCTION

Software visualization is a well-explored area of research which has been applied to several aspects of computing. Visualization has been shown to improve the productivity and effectiveness of programmers, especially when working with complex systems which can span the work of several people over extended periods of time [1,9,11]. The object-oriented languages found in such systems can easily obscure the original solution they were used to implement. With such features as polymorphism and dynamic typing, it is important to consider both the compile-time and runtime characteristics of these languages during visualization [14].

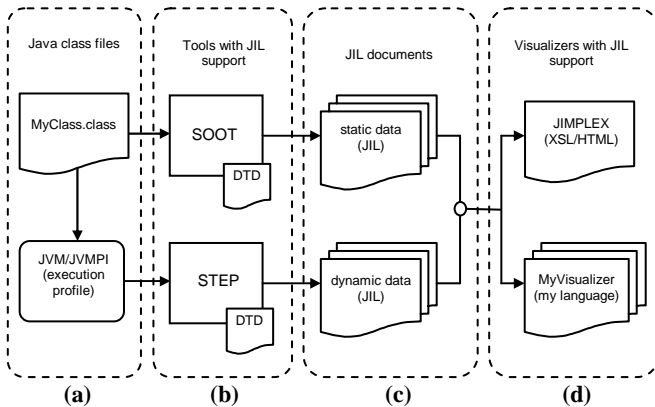
The framework presented in this paper permits the visualization of intermediate representations of Java used by an optimizing compiler. In this environment, the scope of the visualization covers both the software and the underlying characteristics of the execution platform. This data can be extracted from different intermediate languages and other representations, as well as sources of runtime data.

Specifically, we present the following:

- A common intermediate language called *JIL*, capable of encapsulating Java intermediate representations and associating both static and dynamic data with individual code elements.
- Extensions to existing software tools which allow static characteristics and analysis results, as well as dynamic runtime data, to be exported as *JIL* documents.
- A new visualization implementation using *JIL* as a data source, allowing data from the multiple tools to be combined in a customizable interface.

## 2. VISUALIZATION FRAMEWORK

In Figure 1 we present a visualization framework which is designed to be customizable and scalable. The foundation of this system is the new Java Intermediate Language (*JIL*) which can encapsulate existing intermediate languages and related program characteristics. In the following sections we discuss the design of this language, and present two tools which support it and provide suitable sources of both static and dynamic characteristics of the code. We then present an example visualization interface which combines these characteristics by using *JIL* as a data source. Finally, we discuss some of the data management details and present our conclusions.



**Figure 1: Overview of the visualization framework.** (a) Java class files compiled by a Java compiler and executed by a virtual machine. (b) Tools which are able to extract data from Java class files at compile-time and runtime; the data provided by each tool is specified in an accompanying DTD. (c) *JIL* documents containing data extracted in b. (d) Visualizers which use the *JIL* documents as sources of data.

## 2.1 Intermediate Language Design

Intermediate languages are used by optimizing compilers to provide separate modules a representation to work with which is independent of the code being generated. This encourages modularity throughout the compiler, and allows the code generator to be retargeted towards another platform without having to re-implement any analyses or optimizations. Java itself has been described as an ideal intermediate language, providing strong types and other language features which help debug a new language compiler [10].

### 2.1.1 Java Intermediate Language

In order to combine the information which can be associated with these low-level languages, a common format was required which could act as metadata for code. We designed *JIL* for the purpose of encapsulating intermediate representations of Java source code. Although strictly defined, *JIL* remains extensible and able to support undiscovered extensions. Its use differs from traditional intermediate languages in that its content is independent of the tools which use it. *JIL* is not designed to be manipulated or decompiled into bytecode, but provides a bridge between visualization interfaces and tools which use traditional intermediate languages. We only describe the relevant highlights of *JIL* in this document, however a complete specification is available separately as a technical report [8].

*JIL* is based on the Extensible Markup Language (XML) and benefits from many of the features of this well established format, including a growing number of tools and APIs as well as support coming from a large community of developers [2]. Like XML, *JIL* documents are portable across platforms and networks, with native support in most modern web servers and clients. Compatibility is achieved by defining language semantics and restrictions using document type definitions (DTDs). DTDs are simple to extend, and provide a strict method of enforcing compatibility between *JIL* tools; applications which use this standard XML schema

for validating their input can identify which extensions to expect from a *JIL* document.

*JIL* documents are composed of a hierarchy of nested tags which describe the layout of a Java class. This includes enumerations of the fields and methods of the class, including elements which are only found in lower-level intermediate languages, such as labels. This basic skeleton of code elements provides a framework upon which language extensions can be added. These extensions can include both static and dynamic characteristics of the code, and allow such information to be associated directly with each code element.

In order to demonstrate the kinds of extensions supported throughout this paper, we consider polymorphism as a characteristic of the code with both static and dynamic properties. Figure 2a shows a fragment of Java code where polymorphism can be explored at compile-time by analyzing the potential targets of call sites, while the runtime behavior of these sites can reveal which targets are actually being invoked. The following sections continue this example and describe two tools capable of expressing the results of such analyses as *JIL* extensions.

## 2.2 Static Code Elements

All *JIL* documents are required to contain some basic elements of a Java class file. These elements make up the fundamental structure of the document, upon which annotations can be added. Most common code annotations are analysis results which can be calculated by inspecting bytecode. These are static facts and characteristics of the code which are known once the Java source is compiled.

### 2.2.1 Generating static data

In the top portion of Figure 1b we see a code tool capable of performing static analyses on Java bytecode being passed a compiled class file. These results are then exported as extensions in the *JIL* documents it produces. These extensions are defined in an accompanying DTD, allowing tools to validate and identify the *JIL* documents being produced. By associating extensions with the tools that generate them, several tools can annotate the same code elements independently. Redundant or related extensions may also be compared or combined by tools which support them.

As part of the framework presented in this paper, we added support for the output of *JIL* to *SOOT*, an existing optimization framework which uses intermediate languages to perform static analysis and transformations on Java bytecode [15, 16, 18]. Most of these analyses are performed on *Jimple*, an intermediate language where some basic optimizations can be applied before generating bytecode [17]. For example, analysis of which variables are live at each statement can be used to perform copy propagation and simple aggregation, where unneeded variable definitions can be collapsed. With the addition of *JIL* as an output format, *SOOT* provides an ideal source of static data.

Continuing our running example of polymorphism as a code characteristic, *SOOT* supports several types of static analyses which can associate call sites with potential targets. Call sites within the *JIL* documents produced by *SOOT* can be easily identified as monomorphic or potentially polymorphic by extending each site with these analysis results. Figure 2b shows a simplified example of *JIL* describing a call site extended with the results of class hierarchy analysis (CHA) [7] and variable type analysis (VTA) [13].

```
(a) // myclass.java
...
myInterface myObject;
if( branch1 )
    myObject = new A();
else if( branch2 )
    myObject = new B();
else
    myObject = new C();
myObject.myMethod();
...

(b) <!-- JIL: SOOT output of myclass -->
...
<statement>
<jimple>interfaceInvoke 48.myInterface:
    void MyMethod()</jimple>
<soot_invoketargets method="myMethod">
  <targets analysis="CHA" count="3">
    <target class="A"/>
    <target class="B"/>
    <target class="C"/>
  </targets>
  <targets analysis="VTA" count="2">
    <target class="A"/>
    <target class="B"/>
  </targets>
</soot_invoketargets>
...
</statement>
...

(c) <!-- JIL: STEP output for myclass -->
...
<statement>
<jimple>interfaceInvoke 48.myInterface:
    void MyMethod()</jimple>
<step_callsite method="myMethod">
  <targets count="2">
    <target class="A" invokecount="55"/>
    <target class="B" invokecount="45"/>
  </targets>
</step_callsite>
...
</statement>
...
```

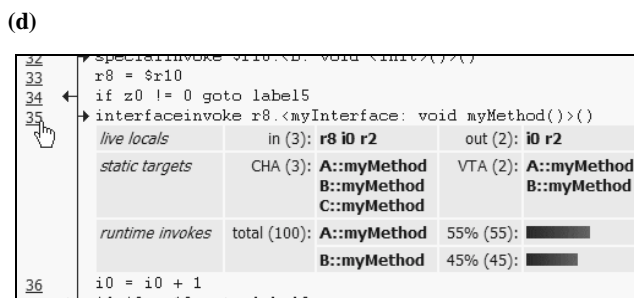


Figure 2: (a) A simple polymorphic Java fragment with an interface invoke; (b) A *JIL* fragment produced by SOOT indicating the possible targets of the call site; (c) A *JIL* fragment produced by a STEP profiling agent indicating the actual number of runtime invokes; (d) A slice from the *JIMPLEX* interface focused on the call site, where the user can browse the extensions from b and c.

## 2.3 Dynamic Code Elements

Code characteristics discovered at runtime provide insight when studying the optimization of object-oriented programs. Languages such as Java include many expensive features, such as polymorphism and garbage collection, which can drastically affect performance. The most effective optimizations target the expensive bytecodes which are generated by Java compilers to provide these features. However, these optimizations can not always be implemented based on static information alone.

### 2.3.1 Collecting dynamic data

The bottom portion of Figure 1b shows runtime data being collected by a profiling tool which supports *JIL* output. Annotation of dynamic data in *JIL* works much like with static data. In order to manage both static and dynamic extensions, each *JIL* document contains a history of all the tools and operations which have authored it. This allows dynamic data from separate executions of the same tool to coexist in a single document. Tools which support *JIL* uniquely identify each execution with an entry in the document's history. These entries typically include the profiling or benchmarking agent which was used, a timestamp, and information about the execution environment. This allows *JIL* tools to compare and process multiple runtime data sets.

In order to collect dynamic data within our framework we extended a tool called *STEP*, a customizable profiling framework for evaluating the performance and behavior of object-oriented applications [3–5]. It helps developers build custom profiling agents which collect data describing the runtime behavior of programs; this allows the rapid development of a variety of profilers which apply to both standard and unconventional profiling tasks.

*STEP* provides its own event-based language and accompanying compiler. Profiling agents define events using this language and pass them into an event pipe where the data is compressed and prepared for consumption. We provided a backend to this event pipe which can generate *JIL* documents. The code elements within these documents are annotated with the runtime data collected by the profiling agents. The *JIL* generator is an easily implemented example of an event pipe consumer, and the extensible nature of the profiling framework is well suited to preserve its output in an extensible document format.

We can now revisit our polymorphism example and use *STEP* to generate *JIL* dynamic data. The *JIL* fragment in Figure 2c demonstrates how *STEP* associates actual runtime invocation targets to a particular call site. Although our static class hierarchy analysis indicated that this call site had three potential targets, according to the data collected from this execution run only two targets were invoked. Runtime data must include some extra information describing the execution environment and any other details that are required to uniquely identify the data associated with it. The separation and management of dynamic data is described in section 3.

## 2.4 Visualization of Static and Dynamic Data

We have demonstrated how *JIL* can combine static and dynamic data, and how existing tools can be extended to produce *JIL* documents. These documents provide versatile data sources which are easily merged and accessed, facilitat-

ing the creation of a variety of visualization interfaces.

### 2.4.1 JIMPLEX

Developers working with *SOOT* and *Jimple* required a tool which would allow them to develop and debug optimizations. Once we had extended *SOOT* and *STEP* to produce *JIL* documents, we developed a visualization implementation called *JIMPLEX* with the goal of providing a customizable visualization framework for browsing *Jimple*. *JIMPLEX* is an XML application which uses XSLT to stylize and transform *JIL* documents, delivering visualization interfaces as HTML to common web browsers [6]. XML transformations can be performed on the fly in modern browsers, allowing the interface to be customized without having to recompile code or learn a complex API. These technologies also encourage collaboration, allowing both the data sources and the visualization interface to be shared between platforms and devices across the Internet.

By using the *JIL* documents produced by *SOOT* and *STEP* as data sources (seen in Figures 2b and 2c), *JIMPLEX* can browse *Jimple* annotated with data from both tools. Figure 2d is a screenshot of the *JIMPLEX* interface, where the user is focused on the *Jimple* statement containing a polymorphic call site. By comparing the results of static analyses to the actual runtime behavior we can verify that the variable type analysis performed by *SOOT* was accurate in eliminating the potential target from class *C*, based on the last profiling run.

Figure 3 is a screenshot of *JIMPLEX* visualizing a Java class in a web browser. Two *JIL* documents are used as data sources to dynamically generate the interface. Static information, such as the relative number of statements per method or label is derived from the *JIL* document produced by *SOOT*. The runtime data, such as the cumulative number of method invokes and field accesses, was extracted from another document produced by a *STEP* profiling agent. By exposing both kinds of data in the same interface the user can identify possible optimization targets, such as the relatively small method `isInteger`, which was executed the most frequently.

## 3. MANAGING JIL DOCUMENTS

The framework presented here encourages interoperability between code tools, regardless of their implementation or supported languages. Any number of tools can contribute to a *JIL* document, allowing collaboration and modularity between tools which would normally be difficult to achieve. Adding support to existing tools requires very little implementation, and many APIs exist which can already validate, parse, and generate compliant *JIL*. The following sections discuss the non-trivial task of providing this kind of interoperability in more detail.

### 3.1 Document Structure

*JIL* documents use the nesting structure of XML to represent the hierarchy of code elements in a Java class. The basic structure of *JIL* documents is illustrated in Figure 4. The current *JIL* definition requires that documents contain a base structure consisting of several required elements. These describe the basic code elements which all classes contain, such as enumerations of fields and methods. By nesting elements, *JIL* can contain optional extensions to these base elements which will not break the underlying structure of

Representation	Lines	File size
Java source	7	140 bytes
<i>Jimple</i> source	23	578 bytes
<i>JIL</i> document	151	7,076 bytes

**Table 1: Relative sizes of representations of a simple Java program.**

the document when removed. This makes document extensions easier to manage, since they can be swapped in or out, and new extensions can be introduced while maintaining backwards compatibility with previous versions of the same document.

*JIL* documents can achieve this kind of manageability by strictly defining the grammar of their contents and self-describing any included extensions. A DTD is used to define the restrictions on which elements and attributes can and must be included. This DTD can also be extended with references to extension definitions, allowing the base *JIL* grammar to evolve independently of any extensions. Definitions are also independently versioned, so that validation can indicate which generation of *JIL* to expect or which extensions are available. Following the trend of other XML technologies, DTDs can be referenced locally or across the internet during validation.

One of the major drawbacks of the XML format is its verbosity. Much additional text is required in order to represent the elements and attributes which define the structure and hierarchies of a *JIL* document. As an example, Table 1 shows the relative sizes of three representations of the same Java program. Although *JIL* documents are significantly larger than the Java source code they represent, they are effectively compressed using standard compression techniques.

### 3.2 Merging

There are several different kinds of merging which are possible when combining data from multiple *JIL* documents. In some cases, an entire package of *JIL* documents might describe a class to be visualized. A single class could also be represented as separate *JIL* documents, and then later combined during processing in order to improve performance or manageability. This section describes the different types of merging which are possible.

The most straightforward merging involves *JIL* documents which represent different classes; in this case there is no merging required. Data in these documents does not intersect or conflict since it is referring to an entirely separate class. Tools that want to browse a complete package or compare the data collected on separate classes can load each document and process them independently. There are no notable caveats in this case, and the details of how to handle each document is left up to the implementation.

Once a tool is handling *JIL* documents which refer to the same class, they can start to take advantage of the object-oriented document model when merging. Such tools typically treat these documents as a single entity describing a Java class. Each document's extensions and data can then be associated with this common entity, and their interactions are left up to the implementation. A simple union of all the data can be performed which presents the user with

## public class Prime extends java.lang.Object

### Methods (9)

declaration	name	statements	calls (622)
public void <init>()	<u>[init]</u>	3	0
public static void main(java.lang.String[])	<u>main</u>	104 ██████████	1
private static java.lang.Long calcPPS(long, long)	<u>calcPPS</u>	31 █████	104 █████
private static boolean isInteger(double)	<u>isInteger</u>	8 █	406 ██████████
public static boolean isPrime(long, int)	<u>isPrime</u>	57 ████████	103 █████
public static void printOut(java.lang.String)	<u>printOut</u>	4	7
public static void printToFile(java.lang.String)	<u>printToFile</u>	3	0
private static void writeFile(java.lang.String, java.lang.String)	<u>writeFile</u>	20 ███	1
private static void fatalError(java.lang.String)	<u>fatalError</u>	5	0

### Fields (4)

modifiers	type	name	reads (247)	writes (911)
private	double	<u>_currentPrime</u>	103 ██████████	509 ██████████
private	long	<u>_currentPPS</u>	107 ██████████	209 ██████████
private	long	<u>_maxPPS</u>	30 █████	106 █████
private	long	<u>_minPPS</u>	7 █	80 ███

### :::main

public static void main(java.lang.String[]) [top](#)

### Statements (104)

```

0  r0 := @parameter0: java.lang.String
1  $r1 = new java.lang.Long;
2  specialinvoke $r1.<java.lang.Long:
   void <init>(long)>(100L);
3  r2 = $r1;
4  $i5 = lengthof r0
5  ← if $i5 == 1 goto
6  staticinvoke <Prime: void printOut
   (java.lang.String)>(" usage: java
   prog <# of primes to find>");
7  staticinvoke <Prime: void printOut
   (java.lang.String)>(" defaulting to
   100 primes...");
8  ← goto label1;
9  $r6 = new java.lang.Long;
10 $r7 = r0[0];
11 specialinvoke $r6.<java.lang.Long:
   void <init>(java.lang.String)>($r7);

```

### Labels (10)

name	statements
default	9 █████
label0	4 █
label1	12 █████
label2	12 █████
label3	6 █
label4	10 █████
label5	6 █
label6	1
label7	17 █████
label8	27 ██████████

### Parameters (1)

### Locals (62)

Figure 3: Screenshot of *JIMPLEX* browsing an intermediate representation of a Java class; the interface is generated dynamically from separate *JIL* data sources. Enumerations of the fields and methods of the class include runtime data, such as method invokes and field accesses. The statement listings for each method expose static analysis results to the user, such as variable liveness and flow information.

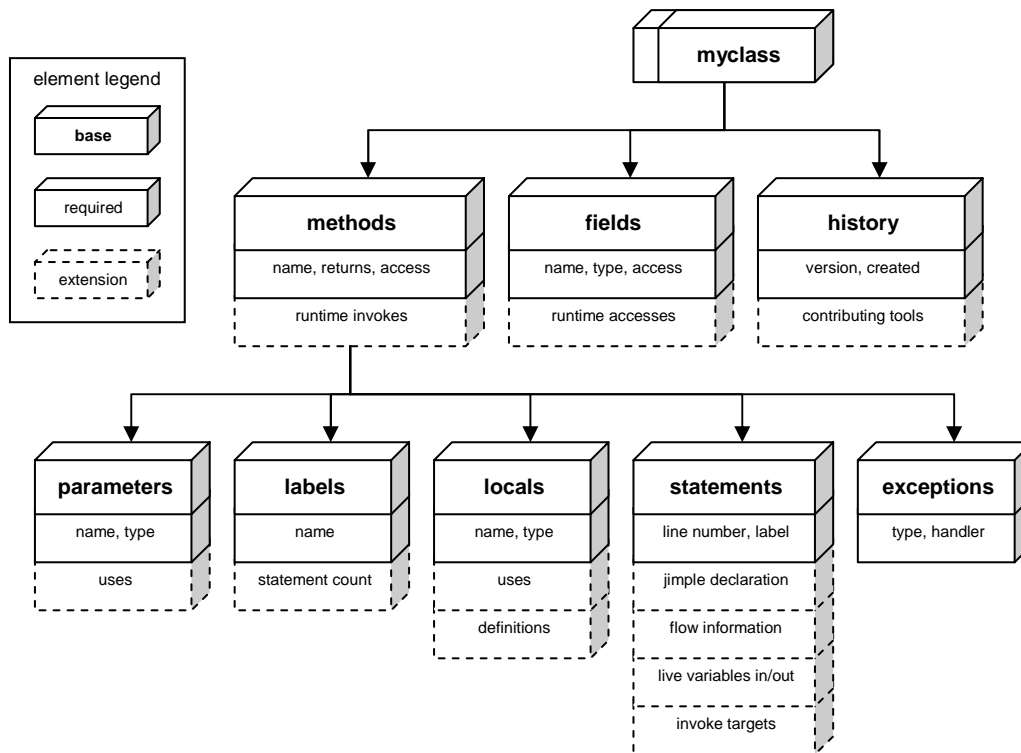


Figure 4: *JIL* document structure. The single inheritance hierarchy of XML is ideal for modelling Java constructs. Base elements provide a structure upon which intermediate languages and metadata can be added. Required elements contain basic information about base elements which is necessary for identification. Language extensions can be associated with any element and provided by arbitrary tools.

a single class representation which includes all the extensions from each document. This is a common case when a tool is passed *JIL* documents from separate sources, each containing different extensions on the same class. These extensions can be combined when loaded by the tool in order to hide their logical separation from the user. This can be convenient when visualizing multiple documents from different remote sources, where their physical separation becomes more of a convenience. For example, if one tool is still being developed and debugged, its extensions can be kept separate from those generated by other more stable tools. The ability to separate extensions in this way is also convenient for research groups working on different extensions independently across the Internet.

In some cases tools can generate document extensions which intersect, meaning they describe the same characteristics of the code with different empirical data. Such data is typically collected at runtime by profiling or benchmark tools over several execution runs, while visualization is targeted at the same code objects. Visualizers can display averages and other calculations by identifying and processing this intersecting data. This process requires some basic algorithms used by the visualizer to decide how to combine the data and present the user with code characteristics of interest.

The visualization framework presented in this paper separates the interface from the data. An open data represen-

tation allows custom interfaces to define how the user visualizes intersecting data. The format and structure of *JIL* is designed to give interfaces more flexibility when deciding how to interpret the data. Simple interfaces can allow basic filtering of datasets where the user can browse the evolution of the code's performance, while complex interfaces might use statistical operations and graphics in order to provide a more comprehensive representation. *JIL* tools which can offer some insight on the interpretation of multiple *JIL* documents can export any data they produce as additional *JIL* extensions. For example, given a *JIL* document describing the local variables which are live at each statement, another tool could interpret this data and export an additional *JIL* document containing lists of variables which are unused and could be eliminated in each method. By chaining the processing and interpretation of *JIL* documents, visualizations can become more complex and cover a larger scope of code characteristics. This also allows a many-to-one relationship between code tools and visualizers.

The process of merging *JIL* document extensions is not trivial, but is facilitated by the wide array of libraries and APIs which can process and parse XML. Many basic combinatory operations are supported by basic interface languages such as XSLT [6] and PHP [12]. These kinds of interpreted languages can allow quick prototyping of new and experimental visualizations. By using a scalable data format, *JIL* tools can process as many documents or extensions as re-

quired by the visualization. Most APIs also support the loading of documents using the well established HTTP protocol for network transmission. This encourages visualizers to support the visualization of remote *JIL* documents, allowing collaboration between tools which might exist on different machines or networks.

The labelling and versioning of *JIL* document extensions allows tools to identify and perform advanced operations when merging and combining data. The following section describes the details of how the current implementation of *JIL* handles the versioning of code elements and extensions.

### 3.3 Versioning

As unambiguous descriptions of Java classes, *JIL* documents allow tools to construct a hierarchical structure of code elements. Document type definitions allow the format of these elements to be recognized and validated using existing XML parsers. DTDs can be referenced using a Uniform Resource Locator (URL), allowing *JIL* tools to provide a unique specification for the *JIL* they support online. Versioning of DTDs allows tools to identify documents based on different versions of *JIL*. Extensions are versioned independently of each other, allowing *JIL* documents to be formed from any combination of supported extensions.

Each *JIL* document contains a history of contributors. This history is a list of tools with attributes which uniquely identify a set of tags within the document. A *JIL* tool uses the document history to describe the operation or command it performed when generating the tags in the document. Version information is usually associated with a history element, which indirectly represents the output of a particular version of a tool. This allows code elements and extensions to be traced back to a specific tool, and then separated by a visualizer when parsing the document. By maintaining this versioned history within each *JIL* document, it allows tools to manage and separate both supported and unsupported extensions.

## 4. CONCLUSIONS

We have presented an open framework for developing visualization and software understanding interfaces. The key features of the framework allow existing and future tools to contribute both static and dynamic code elements to these visualizations, allowing interfaces to be developed based on the information the user wants to analyze rather than what information is immediately available. Collaboration and interoperability are facilitated by using an extensible and portable document format for storing and separating data. This format also allows a many-to-many relationship between applications which can generate and consume *JIL*; several applications can contribute data to a *JIL* representation of a Java class, and several visualizers can then interpret this *JIL* document and present separate interfaces to users.

This framework has been applied to both a static analysis and profiling tools in order to combine the data they provide into a portable, customizable visualization interface. With the addition of *JIL* support, these tools have benefited from the addition of a visualization backend and the ability to export and preserve the code characteristics they can extract. The visualization interface presented in this paper has demonstrated the interoperability and customization that is possible when using *JIL* as a data source, as well

as the usefulness of exposing both static and dynamic data to the user.

### 4.1 Current Progress

Current visualization tools are based on the *JIL* specification 1.0 available online as a Document Type Definition at: <http://www.sable.mcgill.ca/jil>. The *JIMPLEX* visualization framework is also available at this URL as a package of client and server-side scripts to provide visualization interfaces supported by common web browsers. The current release of *SOOT*, which supports *JIL* as an output format, is available at: <http://www.sable.mcgill.ca/soot>. Documentation and release information for *STEP* is available at: <http://www.sable.mcgill.ca/step>.

### 4.2 Future Work

As an open framework, there are many different areas for future work. The *JIL* specification itself is in its infancy, and although only basic extensions are included, the specification is designed to be extended based on the tools that support it. Current support is limited to *SOOT* and *STEP*, but any tool which can provide some insight into software understanding can extend the *JIL* definition and generate *JIL* data. Both static and dynamic extensions are easily specified by DTDs or another form of XML schema. When adding support to a tool for generating or modifying *JIL* documents with data extensions, supplying a DTD allows other tools to validate and recognize those extensions.

Let us consider a user who wants to inspect some generated code in relation to its benchmarking data. They should first decide what code elements would be associated to their benchmarking results, such as extending individual methods with timing information. These extensions should be defined in a DTD, and support for *JIL* should be added to their benchmarking suite using a popular XML API in their favorite programming language. The extension DTD should then be used by a visualization interface to validate *JIL* documents containing these dynamic extensions. The visualizer could present the user with statistical information based on the benchmarking data recorded in the *JIL* documents. Future improvements to the code generator could then be evaluated by comparing successive *JIL* documents containing the benchmarking extensions.

Visualization is the current focus of this framework, however it is only one example of an application for *JIL*. Future use of *JIL* could target any application where meta-data is associated with code. A *JIL*-aware Integrated Development Environment (IDE) might remind the user about methods which are called frequently or suggest the most effective strategy to modularize the code. Software development rarely involves the inspection of static analyses beyond compiler errors, and dynamic information is typically not available until changes to the code may be too costly. Much effort is spent debugging software during development, and most developers target a single problem or area of the code when debugging. A debugger which supported *JIL* could preserve such information with the code allowing the developer to reference this data without having to execute another costly debugging run. By combining static and dynamic data into an extensible document format, tools can provide a developer with information and insight normally obscured by the code.

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